

COEN-6501 DIGITAL SYSTEMS AND SYNTHESIS

Final Examination, Dec. 13, 2012

Examiner: A. J. Al-Khalili

Official Calculators are allowed

Time Allowed: 3hrs

Answer all questions

All questions carry equal weight

Question 1

The following VHDL model illustrates embedded code for generate statements to describe a two dimensional structure.

- Correct all syntax and semantic errors.
- Draw the logic diagram described by the VHDL model.
- Generate the Boolean expression for $X = F(A,B)$.

```
L1  Entity Logic is;
L2  port (A,B: in BIT (3 downto 0) X: out bit);
L3  end LOGIC
L4  architecture STRUCTURE of LOGIC is
L5  component AND2-GATE
L7  port
L6  (A,B: in BIT; Z: out BIT); end component;
L7  component OR_2GATE port (A,B: in BIT; Z out out bit); end component
L8  OR2_GATE;
L9  signal ASIG , OSIG: BIT_VECTOR( 3 downto 0) := X"0";
L10 begin
L11  R: for COL in 1 to 2 generate
L12      C: for ROW in 3 downto 0 generate
L13          R1: if (COL=1) generate
L14              AX: AND2_GATE portmap (A(ROW), B(ROW), ASIG(ROW));
L15          end generate R1;
L16          R1C: if (COL=2 and ROW/=0) generate
L17              OX: OR_2GATE portmap(ASIG(ROW), O(ROW), OSIG(ROW-1));
L18          end generate R1C;
L19          R1C0: if (COL=2 and ROW=0) generate
L20              O0: OR_2GATE portmap (ASIG(ROW), OSIG(ROW),X);
L21          end generate R1C0;
L22      end generate C;
L23  end generate R;
L24  end STRUCTURE;
```

Question 2

The circuit shown in Fig. 1 given below is a Controller. The timing characteristic of the gates and the flip-flops are listed below. The circuit operates at a supply voltage of $5V \pm 10\%$ in an ambient temperature range of 0 to 25 °C with a power dissipation of 1.0 W. The thermal resistance of the package is 40 °C/W and the process variation factor is $\pm 20\%$. Determine:

- All the paths in the circuit.
- The critical path in the circuit. What will be the typical delay of the path?
- Maximum speed of operation for the worst conditions.

Notes:

- Arrival times of inputs: Start, Zero, External is at $-\infty$ and all are registered inputs.
- Temperature degrading factor $M=1.5$
- $K1, K2$ are input and output degrading factors.

Component	T_p (ns)	Input Loading (UL)	$K1$ (ns/UL)	$K2$ (ns/fo)
Inverter	0.15	1	0.08	0.1
NAND	0.3	1.5	0.12	0.15
AND	0.8	2.5	0.12	0.15
D-FF ($t_{su}=0.5$, $t_h=0.2$ ns)	0.7	2	0.10	0.2

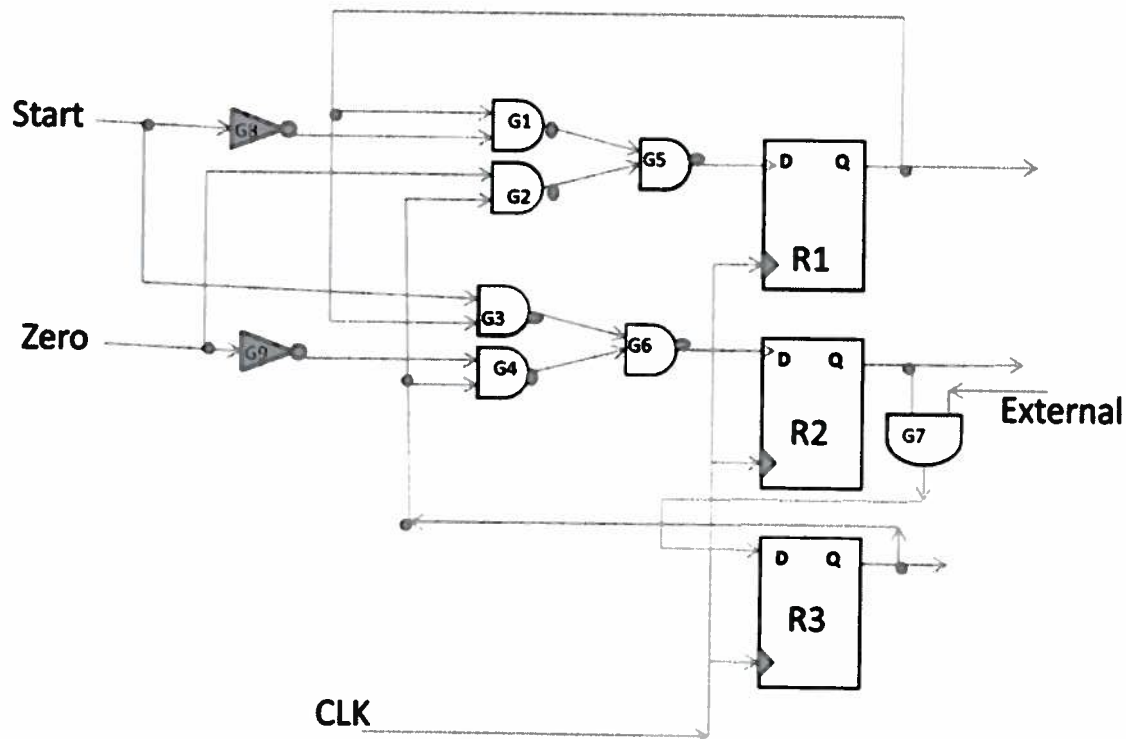


Fig. 1 of Question 2

Question 3

- Design a full Subtractor
- Implement the subtractor in an FPGA with 2 input Look-Up Tables. Implement the same with a 3 input Look-Up Tables. Compare the two results

Question 4

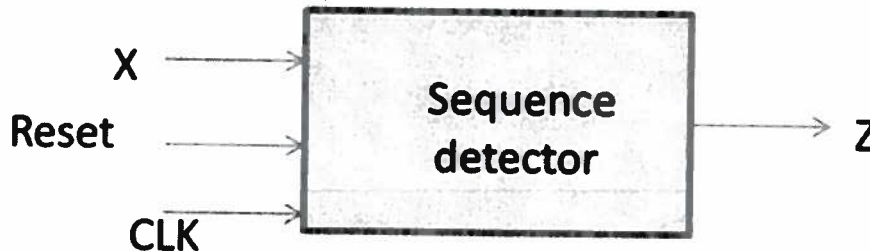
It is required to design a multi-operand adder based on Carry Save Adder (CSA) architecture. Assume the number of operands is equal to five (v,w,x,y,z), each with a width of 4 bits as shown below.

- Construct the adder using a Manchester carry adder for the output stage.
- Evaluate the area in terms of equivalent Full Adder, A and delay in terms FA delay D. assume a MUX is 0.5A, 0.5D and any gate is 0.25A, 0.25D
- Insert a pipeline before the Manchester Carry Adder. Perform the necessary analysis to show the impact of addition of the pipeline register on the speed and area. Assume the following parameters for each Flip Flop in the register: Clock to output delay, $T_{CQ} = 1.5D$, Set up time $T_{su} = D$ and Hold time, $T_h = 0.5D$.



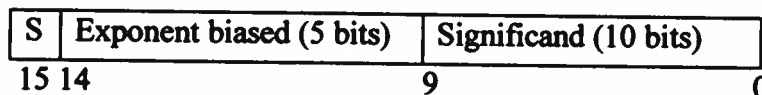
Question 5

Design a Finite State Machine (FSM) to be used as start of a message detector. Serial data, one bit at a time, in packets of 4 bits entering the FSM on line x, are synchronized with the clock and analyzed. The start of data is indicated with arrival of either "1010" or "0101" where an output, z=1 is generated. A separate reset mechanism will put the machine in the initial state.



Question 6

- Represent $A = 0.25$ and $B = -10$ in a floating point format given below assuming a hidden '1'



- Add, $A + B$ in a floating point format and give your result in packaged format.
- Explain the IEEE, rounding method of "to nearest even"
- Give an architecture for floating point Adder.

Question I

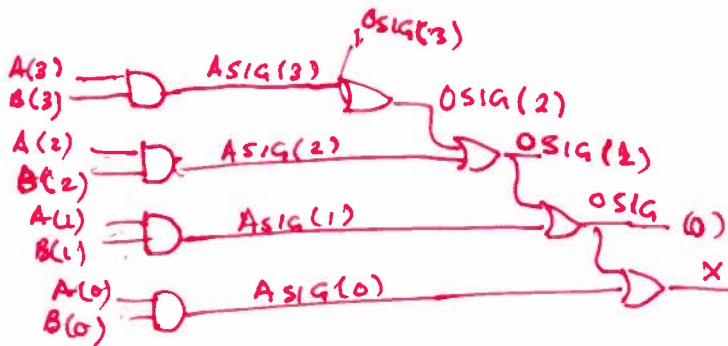
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i) 14)

```

1 L1 Entity Logic is;
2 L2 port (A,B: in BIT (3 downto 0) X: out bit);
3 L3 end LOGIC
4 L4 architecture STRUCTURE of LOGIC is
5 L5 component AND2-GATE
6 L7 port
7 L6 (A,B: in BIT; Z: out BIT); end component;
8 L7 component OR_2GATE port (A,B: in BIT; Z out out bit); end component
9 L8 OR2_GATE;
10 L9 signal ASIG , OSIG: BIT_VECTOR( 3 downto 0) := X"0";
11 L10 begin
12 L11 R: for COL in 1 to 2 generate
13 L12 C: for ROW in 3 downto 0 generate
14 L13 R1: if (COL=1) generate
15 L14 AX: AND2_GATE portmap (A(ROW), B(ROW), ASIG(ROW));
16 L15 end generate R1;
17 L16 R1C: if (COL=2 and ROW/=0) generate
18 L17 OX: OR_2GATE portmap(ASIG(ROW), O(ROW), OSIG(ROW-1));
19 L18 end generate R1C;
20 L19 R1C0: if (COL=2 and ROW=0) generate
21 L20 O0: OR_2GATE portmap (ASIG(ROW), OSIG(ROW),X);
22 L21 end generate R1C0;
23 L22 end generate C;
24 L23 end generate R;
25 L24 end STRUCTURE;
  
```

ii) (g) Component $A(i)$ $B(i)$ $ASIG(i)$ $OSIG(i)$ $OSIG(i-1)$

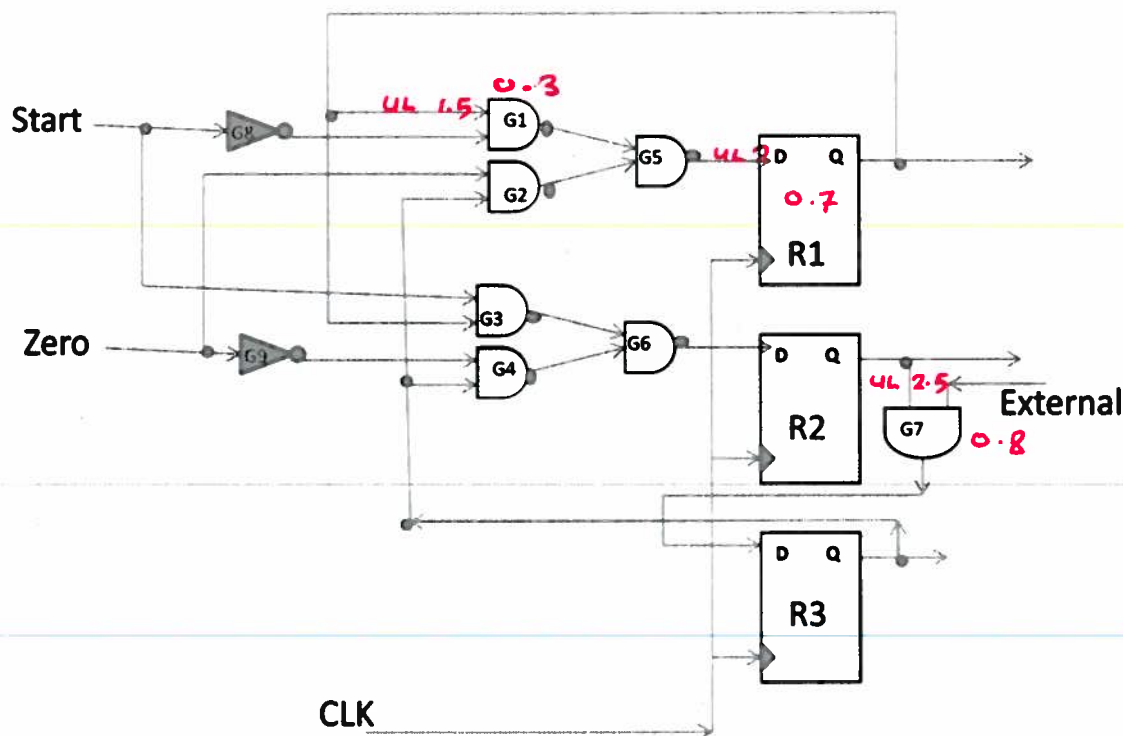


iii) (3) $X = ((A(3) \cdot B(3) + OSIG(3)) + A(2) \cdot B(2) + A(1) \cdot B(1)) + A(0) \cdot B(0)$

Q2

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d) ③

There are 5 paths

- ① $R1 \rightarrow G1 \rightarrow G5 \rightarrow R1$
- ② $R1 \rightarrow G3 \rightarrow G6 \rightarrow R2$
- ③ $R2 \rightarrow G4 \rightarrow G6 \rightarrow R2$
- ④ $R3 \rightarrow G2 \rightarrow G5 \rightarrow R1$
- ⑤ $R2 \rightarrow G7 \rightarrow R3$

Paths ① ② ③ ④ are EQUAL

b) ④ Path ①

$$\begin{aligned}
 &0.7 + 0.2(3) + 2 * 0.1(1.5) \\
 &0.3 + 0.15(1) + 0.12(1.5) \\
 &0.3 + 0.15(1) + 0.12(2)
 \end{aligned}$$

$$+ 0.5 = 3.42 \text{ ns}$$

Path ⑤

$$\begin{aligned}
 &0.7 + 0.2(1) + 0.1(2.5) \\
 &0.8 + 0.15(1) + 0.12(2)
 \end{aligned}$$

$$+ 0.5 = 2.84 \text{ ns}$$

Critical path is path 1. With max frequency of $\frac{1}{3.42} \Rightarrow \underline{\underline{292.4 \text{ MHz}}}$

e) ③ Worst case

$$T_a = 25^\circ\text{C} \quad \theta = 40^\circ\text{C/W} \quad P = 1\text{W}$$

$$P = \frac{T_J - T_a}{\theta} \quad T_J = 65^\circ\text{C}$$

$$f_t = \left(\frac{T_J}{T_a} \right)^{-1.5} \quad \left(\frac{338}{298} \right)^{-1.5} = 1.20$$

Voltage

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$$K_v = \frac{1}{(1 \pm 0.01 \% V_d)}$$

with $\pm 10\%$ worst case is 90% .

$$= \frac{1}{1 - 0.1} = \frac{1}{0.9} = 1.11$$

Process

$$K_p = (1 + 0.01 K_p)$$

with $\pm 20\%$

$$= 1 + 0.2 = 1.21$$

$$K = K_p * K_v * K_T = 1.11 * 1.21 * 1.20 \\ \approx 1.6$$

Worst Case frequency $\approx \underline{\underline{182.75}}$

Q3

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a) ④

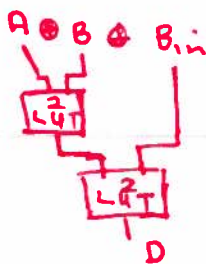
	A	B	B_{in}	D	B_{out}
0	0	0	0	0	0
1	0	0	1	1	1
2	0	1	0	1	1
3	0	1	1	0	1
4	1	0	0	1	0
5	1	0	1	0	0
6	1	1	0	0	0
7	1	1	1	1	1

Two input table

B_{in}	00	01	11	10
0	0	1	4	4
1	1	3	7	5

$$D = A \oplus B \oplus B_{in}$$

Two input table



b) ③

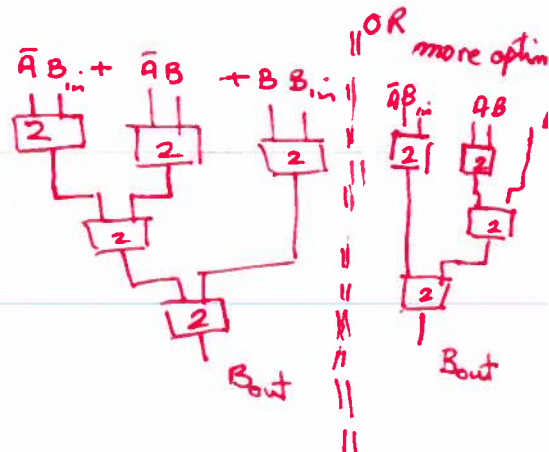
Three input Table



B_{in}	00	01	11	10
0	0	1	4	4
1	1	3	7	5

$$B_{out} = \bar{A}B_{in} + \bar{A}B + B B_{in}$$

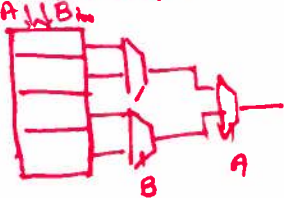
$$\text{Or } B_{out} = \bar{A}B + B_{in}(A \oplus B)$$



Implementation

c) ③

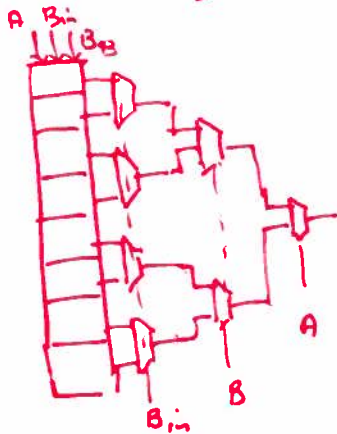
2 input LUT



A2
Area of 2 Variable LUT:
4 Cell + 3 MUX

D2
Delay Read:

1 Cell + 2 MUX



A3
Area of 3 Variable LUT:

8 Cell + 7 MUX

D3
Delay Read:

1 Cell + 3 MUX

Area of 3 LUT

8 Cell + 7 MUX

Delay of 3 LUT

1 Cell + 3 MUX

Comparison

Comparison

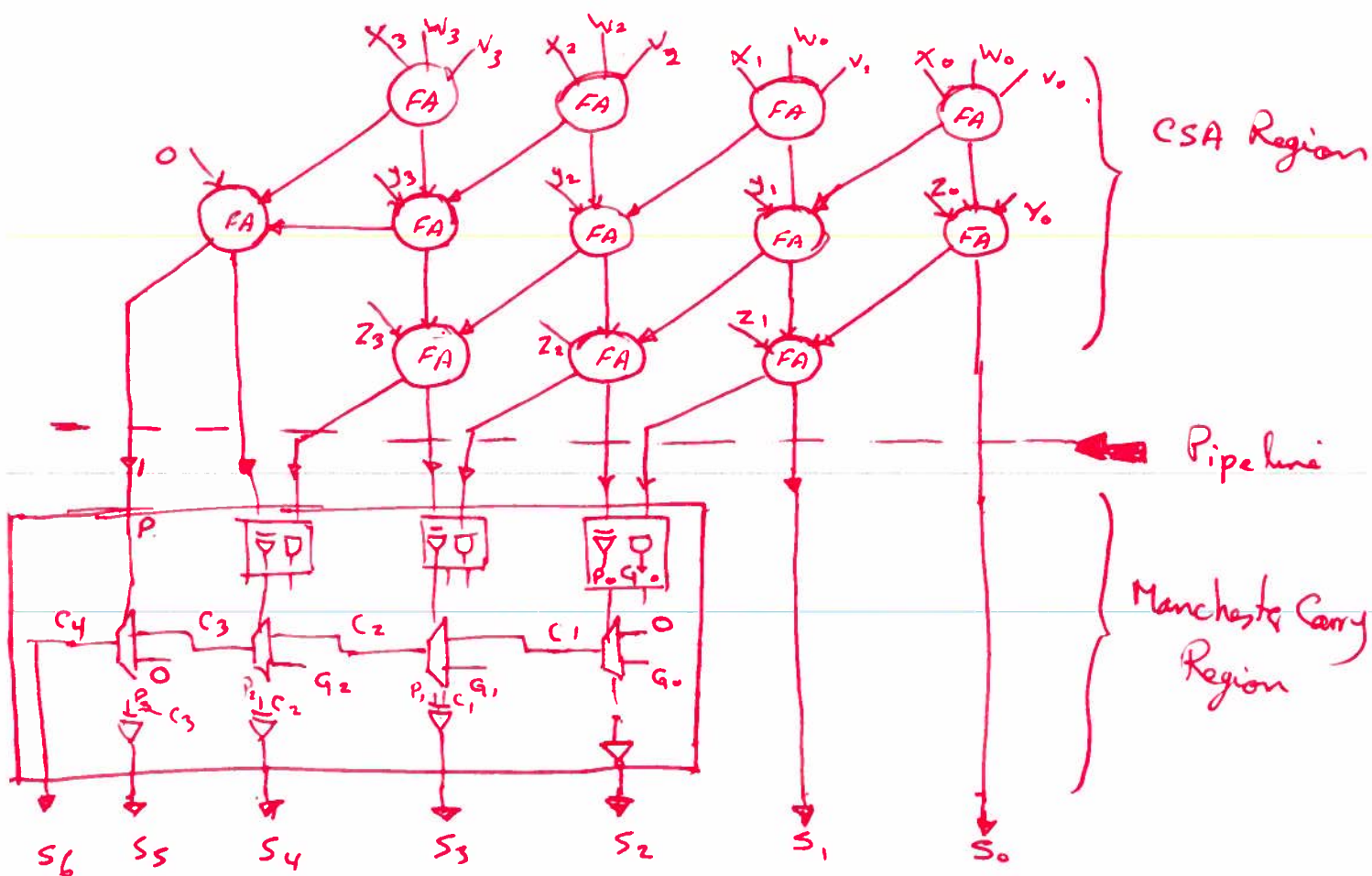
Area of 2 LUT

7 A2 = 28 Cell + 21 MUX

Delay of 2 LUT

1 Cell + 6 MUX

Using 3-input LUT is faster and more saving in area



b) Delay = $3D + 0.25D + 4 \times 0.5D + 0.25 = 5.5D$

CSA PG MUX₀ XOR

Area = $12A + 10 \times 0.25A + 5 \times 0.5A = 17A$

FA Gates MUX₀

c) When we insert the pipeline, the Combinational logic is broken into 2 parts

First part will be the CSA = $3D$

Second " " " " Manchester = $4 \times 0.5D + 0.25D + 0.25D = 2.5D$

Second part is less than the first part delaywise so

$T = 1.5D + 3D + D = 5.5$

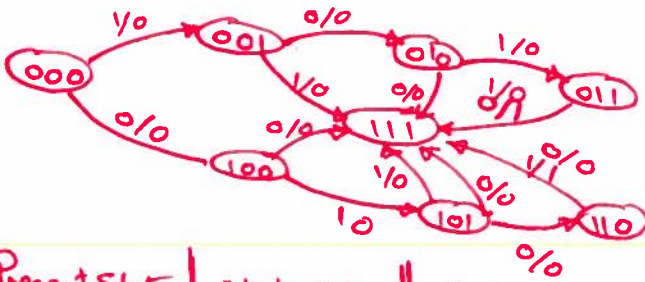
T_{eq} CSA t_{su}

Delay is not improved yet, the area is increased by 9 FF.

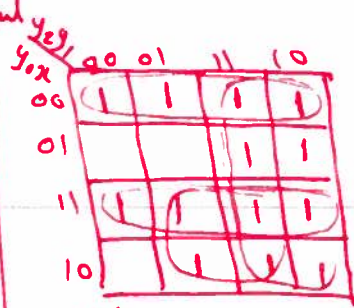
Q5

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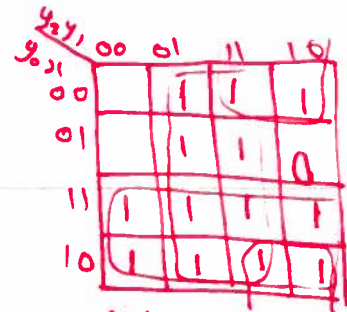
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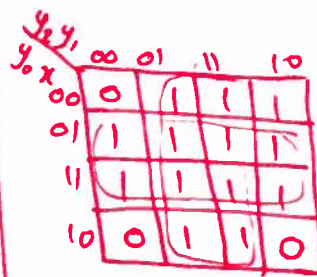
min term	Present State				Next State			output Z
	y_2	y_1	y_0	x	y_2^+	y_1^+	y_0^+	
0	0	0	0	0				
1	0	0	0	1				
2	0	0	1	0				
3	0	0	1	1				
4	0	1	0	0				
5	0	1	0	1				
6	0	1	1	0				
7	0	1	1	1				
8	1	0	0	0				
9	1	0	0	1				
10	1	0	1	0				
11	1	0	1	1				
12	1	1	0	0				
13	1	1	0	1				
14	1	1	1	0				
15	1	1	1	1				



$$y_2^+ = y_2 + \bar{y}_0 x + y_0 x + y_1 y_0$$

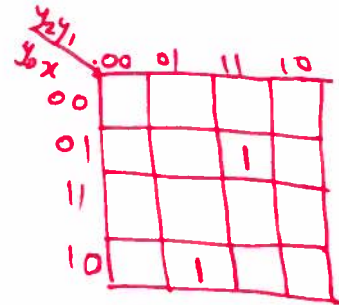


$$y_1^+ = y_1 + y_0 + y_2 \bar{x}$$

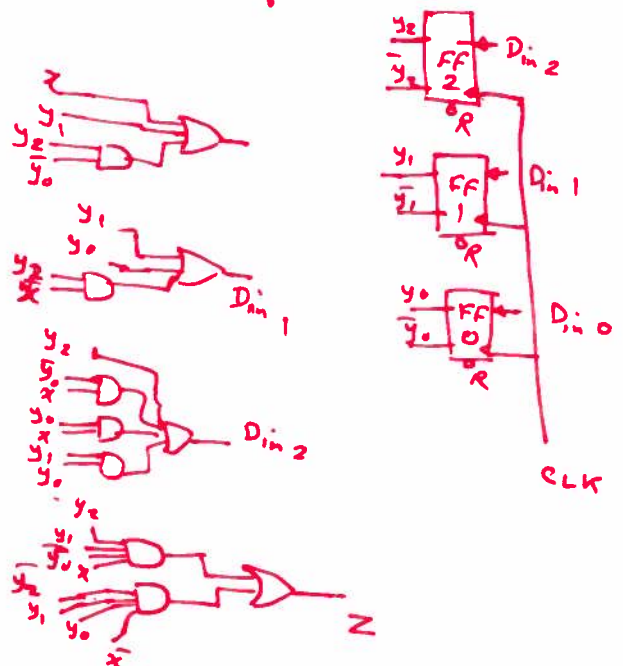


$$y_0^+ = y_1 + x + y_2 \bar{y}_0$$

Using D Flip Flop $y_i^+ = D_{in i}$



$$Z = y_2 y_1 \bar{y}_0 x + \bar{y}_2 y_1 y_0 \bar{x}$$



a) ③ Bias = $2^{n-1} - 1 = 2^{5-1} - 1 = 15$

$A = 0.25_{10} = 0.01_2 = 1.0 \times 2^{-2}$

$B = 10_{10} = 1010.0_2 = 1.010 \times 2^3$

$e_A = 15 - 2 = 13_{10} = 1101_2$

$e_B = 15 + 3 = 18_{10} = 10010_2$

Representation:

A →

0	0	1	1	0	1	0	0	0	0	0	0	0	0	0	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

B →

1	1	0	0	1	0	0	1	0	0	0	0	0	0	0	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

b) ③

A+B exp2 - exp1 = 3 - (-2) = 5 Shift A 5 places to the right

B is -ve Complement it

01.01000 2' Complement

110.10111
+ 00001 + 1

110.11000

Now add A+B

-ve number

Complement & add 1

110.11000
+ 00001

 110.11001
+ 001.00110

 1.00011

A+B

1	1	0	0	1	0	0	0	1	1	0	0	0	0	0	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

②

⑥) We round the numbers if $R(M_0+S)$ is true. this translate to the following table

Number	rounded	error	Number	rounded	error
X0.00	X0.	0	X1.00	X1	0
X0.01	X0.	-1/4	X1.01	X1	-1/4
X0.10	X0.	-1/2	X1.10	X1+1	+1/2
X0.11	X1	+1/4	X1.11	X1+1	+1/4

This gives zero error assuming there are equal even and odd numbers

d) ② check notes:

