

## Arithmetic Operations

- We will review the arithmetic building blocks we have previously used, and look at some new ones.
  - Addition
  - incrementer
  - Addition/subtraction
  - decrementer
  - Comparison

BR 8/99

---

---

---

---

---

---

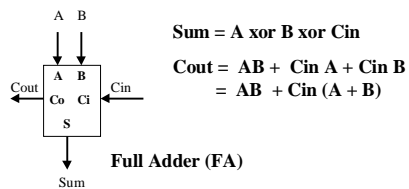
---

---

## Binary Adder

$$F(A,B,C) = A \text{ xor } B \text{ xor } C \quad G = AB + AC + BC$$

These equations look familiar. These define a *Binary Full Adder* :



BR 8/99

---

---

---

---

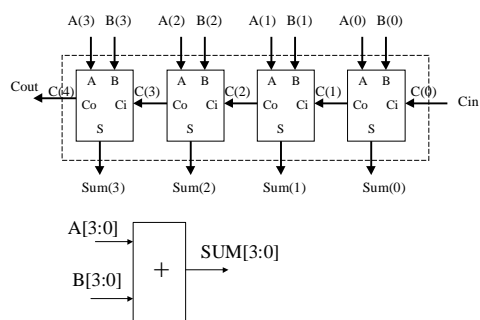
---

---

---

---

## 4 Bit Ripple Carry Adder



BR 8/99

---

---

---

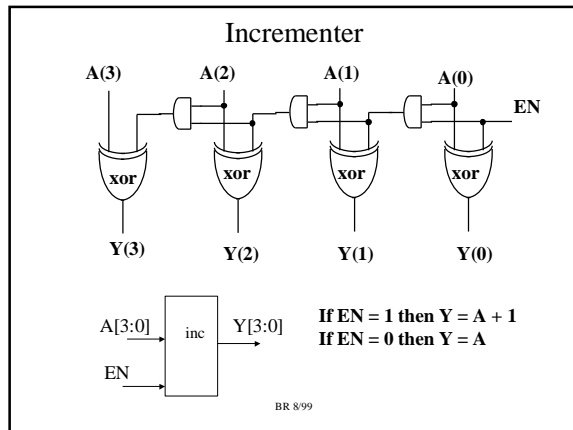
---

---

---

---

---




---

---

---

---

---

---

---

---

### How did we get the Incrementer equations?

**Full Adder equations:**

**Sum** = A xor B xor Cin

**Cout** = AB or Cin A or Cin B  
= AB or Cin (A or B)

Let B = 0, Cin = 1 so that Sum = A + 1. Then equations simplify to:  
**SUM** = A xor 1 xor 0 = A xor 1 = A'  
**Cout** = 0 or 1 (A or 0) = A.

If we want an "En" input, then we want SUM = A if En=0, else SUM = A+1 if En = '1'. Filling in the above equations:

**SUM** = A En' or A' En = A xor En  
**Cout** = A En (note that Cout = 0 if En = 0).

The "Cout" of one bit becomes the "En" signal for the next bit!!!!

BR 8/99

---

---

---

---

---

---

---

---

### A Subtractor

What is subtraction?

$A - B = A + (-B)$

How do you take the negative of a number? Depends on the sign representation (signed magnitude, 1s complement, 2s complement). Lets assume 2's complement since it is most common).

$(-B) = B' + 1$

So:

$A - B = A + (-B) = A + B' + 1$

BR 8/99

---

---

---

---

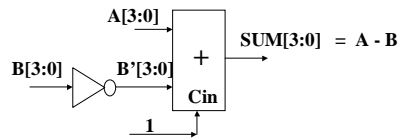
---

---

---

---

### Subtractor using an Adder



What if we want a block that can do both addition and subtraction?

BR 8/99

---

---

---

---

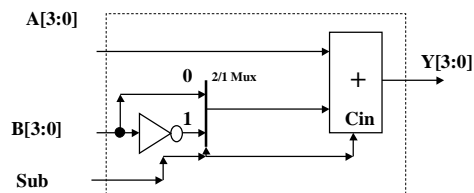
---

---

---

---

### Adder/Subtractor



VHDL representation:

$Y \leftarrow (A-B) \text{ when } (Sub = '1') \text{ else } A+B;$

BR 8/99

---

---

---

---

---

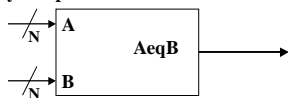
---

---

---

### Recall what a Comparator is...

Equality comparator.



$A=B$  if  $A(0) = B(0)$  and  $A(1) = B(1) \dots$  and  $A(n-1)=B(n-1)$

Recall that "xnor" function is '1' if  $A=0, B=0$  or  $A=1, B=1$ !  
So  $AeqB$  is:

$AeqB = (A(0) \text{ xnor } B(0)) \text{ and } (A(1) \text{ xnor } B(1)) \text{ and } \dots \text{etc.}$

BR 8/99

---

---

---

---

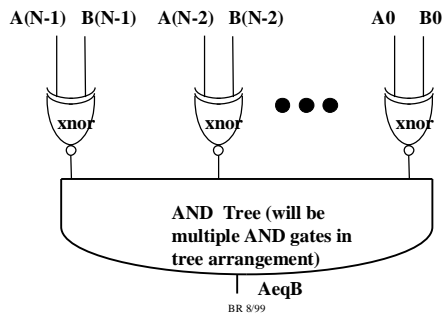
---

---

---

---

What is logic structure for equality comparator?




---

---

---

---

---

---

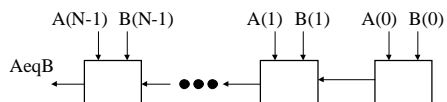
---

---

Is there another Logic structure possible?

Compare "iteratively" from LSB to MSB

If  $A(0) = B(0)$  then  
 if  $A(1) = B(1)$  then  
 ....  
 If  $A(N-1) = B(N-1)$  then  
 AeqB = '1'; !!!!!



Signal from one bit block to next is "enable" for that block.

---

---

---

---

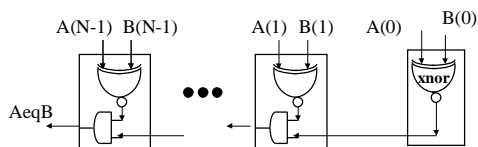
---

---

---

---

Iterative Comparator Structure



An advantage to this structure is that the design for each bit is the same same, and we can extend it indefinitely. But it will be slow.

---

---

---

---

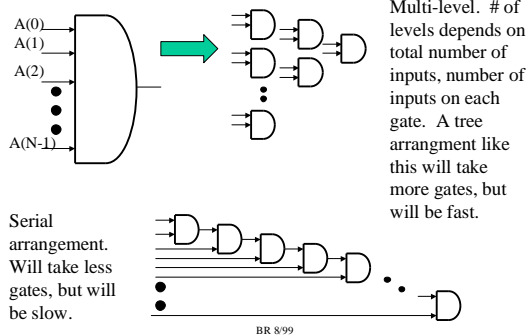
---

---

---

---

## Two ways to do a Large AND function




---

---

---

---

---

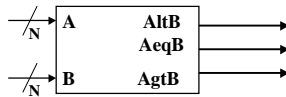
---

---

---

## What about “<” (less than), “>” (greater than?)

### Full comparator.



The logic for AltB, AgtB depends on whether we are comparing signed numbers or not. We will assume unsigned numbers for now.

BR 8/99

---

---

---

---

---

---

---

---

## Logic for “AgtB” (unsigned)

Consider  $A > B$ , both  $N$  bit numbers,  $A[N-1:0]$ ,  $B[N-1:0]$

If  $(A(N-1) = '1' \text{ and } (B(N-1) = '0'))$  then  
 $\text{AgtB} = '1';$   $\leftarrow A=1xxx\dots B=0xxxxx$   
 elsif  $((A(N-1) = B(N-1)) \text{ and } (A(N-2) = '1' \text{ and } (B(N-2) = '0'))$  then  
 $\text{AgtB} = '1';$   $\leftarrow A=01xx\dots B=00xxxx$   
 $\leftarrow A=11xx\dots B=10xxxx$   
 etc...

Look at “bit(i)”. The enable signal from previous bit is  $A = B$  up until now. If this is ‘1’, then we need to do a comparison.

However, if “AgtB” is already true, then we don’t need to do comparison and can skip this comparison!

BR 8/99

---

---

---

---

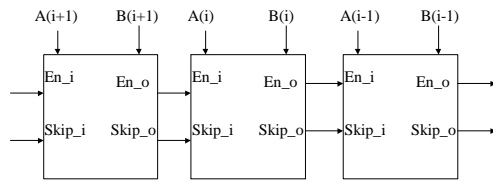
---

---

---

---

## Iterative Implementation of AgtB



$en_o = (A \text{ xnor } B) \text{ and } en_i$ ;  $\leftarrow$  A=B signal

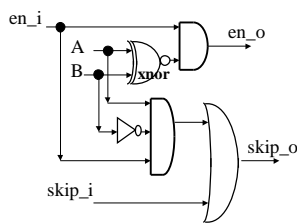
If (skip\_i = '1') then  
 skip\_o = '1';  
 else  
 skip\_o = en\_i and (A and B');  
 end if;

BR 8/99

## Logic Implementation

$en_o = (A \text{ xnor } B) \text{ and } en_i$ ;

If (skip\_i = '1') then  
 skip\_o = '1';  
 else  
 skip\_o = en\_i and (A and B');  
 end if;



BR 8/99

Can use a K-map  
 to simplify this  
 logic.

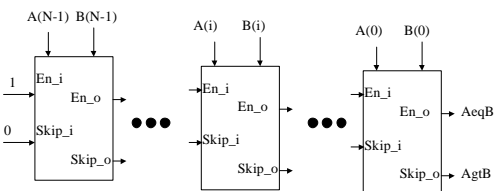
The *skip\_o* of the LAST  
 bit is the AgtB signal!

The *en\_o* of the LAST  
 bit is the AeqB signal!

What about AltB???

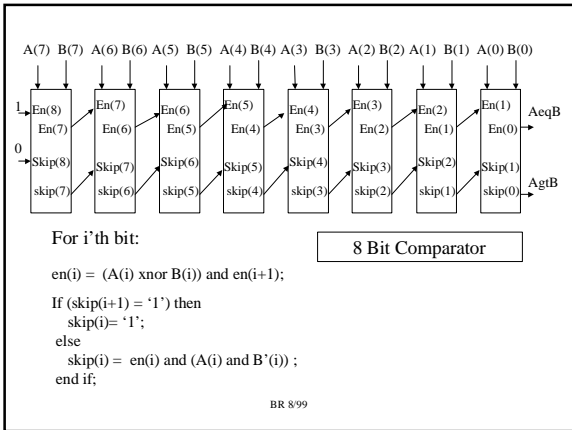
AltB = AgtB' and AeqB'

## Final Comparator



$AltB \leftarrow \text{not } (AeqB) \text{ and } \text{not } (AgtB);$

BR 8/99




---

---

---

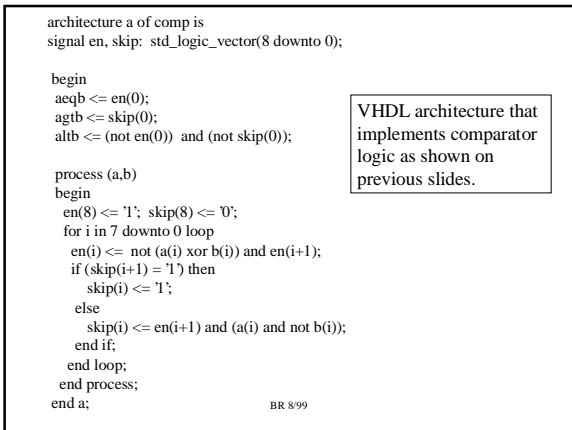
---

---

---

---

---




---

---

---

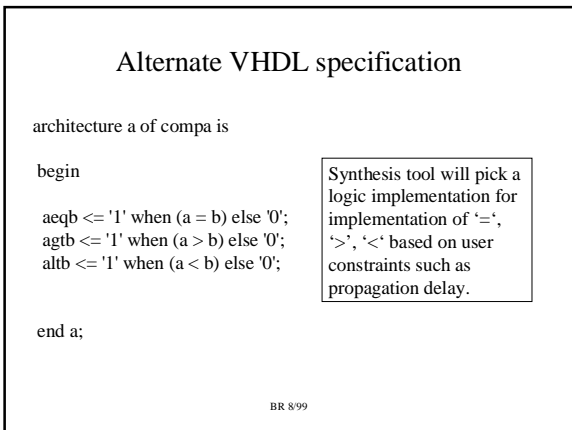
---

---

---

---

---




---

---

---

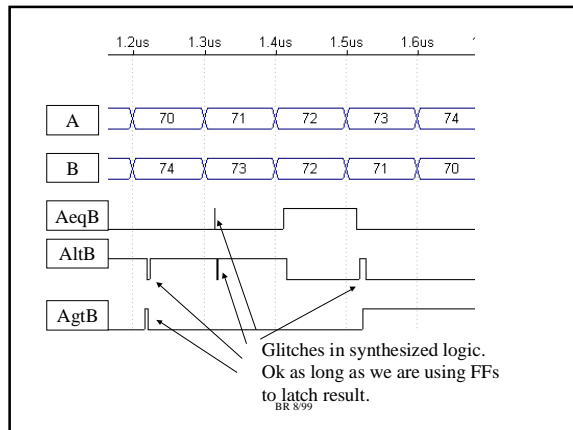
---

---

---

---

---




---

---

---

---

---

---

---